Research Institute for Discrete Mathematics Chip Design Summer Term 2014

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Exercise Set 11

Exercise 11.1:

Consider the LP formulation of the HITCHCOCK TRANSPORTATION PROBLEM

$$\begin{array}{ll} \min & \sum_{i=1}^{n} \sum_{j=1}^{m} c_{ij} \cdot x_{ij} \\ s.t. & \sum_{j=1}^{m} x_{ij} = a_i \\ & \sum_{i=1}^{n} x_{ij} = b_j \\ & x_{ij} \geq 0 \end{array} \qquad \qquad \forall i = 1, ..., n, \\ & \forall i = 1, ..., n, j = 1, ..., m, \end{array}$$

and its dual

$$\max\left\{\sum_{i=1}^{n} a_i \cdot u_i + \sum_{j=1}^{m} b_j \cdot v_j : u_i + v_j \le c_{ij} \,\forall \, (i,j) \in \{1,...,n\} \times \{1,...,m\}\right\}.$$

Prove: Given an optimum dual solution (u, v), an optimum primal solution can be computed in $\mathcal{O}(n)$ time if m is considered to be constant.

Exercise 11.2: Show how the above LP with m = 2 can be solved in $O(n \log n)$ time without calling any min-cost flow algorithm.

Deadline: Thursday, July 3, before the lecture. The websites for lecture and exercises are linked at

http://www.or.uni-bonn.de/lectures/ss14/ss14.html

In case of any questions feel free to contact me at scheifele@or.uni-bonn.de .