Combinatorial Optimization
Winter term 2010/2011

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## Exercises 7

## Exercise 1:

Show that a 2-factor approximation for a (cardinality maximal) b-matching in a graph can be found in linear time.

## Exercise 2:

Let $G$ be a $k$-regular and $(k-1)$-edge-connected graph with an even number of vertices, and let $c: E(G) \rightarrow \mathbb{R}_{+}$. Prove that there exists a perfect matching $M$ in $G$ with $c(M) \geq \frac{1}{k} c(E(G))$.
Hint: Show that $\frac{1}{k} \mathbb{1}$ is in the perfect matching polytope, where $\mathbb{1}$ denotes a vector whose components are all one.
(4 points)

## Exercise 3:

Show that a minimum weight perfect simple 2-matching in an undirected graph $G$ can be found in $O\left(n^{6}\right)$ time.

## Exercise 4:

Let $G$ be a graph, and let

$$
P:=\left\{x \in \mathbb{R}_{+}^{|E(G)|}: \sum_{e \in \delta(v)} x_{e}=1 \quad \text { for all } v \in V(G)\right\}
$$

be the fractional perfect matching polytope of $G$. Prove that the vertices of $P$ are exactly the vectors $x$ with

$$
x_{e}= \begin{cases}\frac{1}{2} & \text { if } e \in E\left(C_{1}\right) \cup \ldots \cup E\left(C_{k}\right) \\ 1 & \text { if } e \in M \\ 0 & \text { otherwise }\end{cases}
$$

where $C_{1}, \ldots, C_{k}$ are vertex-disjoint odd circuits and $M$ is a perfect matching in $G-\left(V\left(C_{1}\right) \cup \ldots \cup V\left(C_{k}\right)\right)$.

Deadline: Tuesday, November 30th, before the lecture.

