## Exercise Set 4

**Exercise 4.1.** Consider the SHORTEST EVEN/ODD PATH PROBLEM: Given a graph G with weights  $c : E(G) \to \mathbb{R}_{\geq 0}$  and  $s, t \in V(G)$ , find an *s*-*t*-path P of even/odd length in G that minimizes  $\sum_{e \in E(P)} c(e)$  among all *s*-*t*-paths of even/odd length in G. Show that both the even and the odd version can be linearly reduced to the MINIMUM WEIGHT PERFECT MATCHING PROBLEM.

(5 points)

**Exercise 4.2.** Consider the MINIMUM COST EDGE COVER PROBLEM: Given a graph G with weights  $c : E(G) \to \mathbb{R}_{\geq 0}$ , find an edge cover  $F \subseteq E(G)$  that minimizes  $\sum_{e \in F} c(e)$ . Show that the MINIMUM COST EDGE COVER PROBLEM can be linearly reduced to the MINIMUM WEIGHT PERFECT MATCHING PROBLEM. (5 points)

**Exercise 4.3.** Let G be a graph with edge weights  $c : E(G) \to \mathbb{R}$  and let M be a matching in G with |M| = k that has minimum weight among all matchings in G that contain exactly k edges. Let P be an M-augmenting path in G with minimum gain. Let  $M' := M \triangle E(P)$ . Prove that M' has minimum weight among all matchings in G that contain exactly k + 1 edges.

(5 points)

**Exercise 4.4.** Let G be an undirected graph with edge weights  $c: E(G) \to \mathbb{R}$ , and let M be a matching so that  $c(N) \leq c(M)$  for all matchings N in G with  $|M| - 1 \leq |N| \leq |M| + 1$ . Prove that then M is a maximum weight matching in G.

(5 points)

**Deadline:** November 7<sup>th</sup>, before the lecture. The websites for lecture and exercises can be found at:

## http://www.or.uni-bonn.de/lectures/ws19/co\_exercises/exercises.html

In case of any questions feel free to contact me at rabenstein@or.uni-bonn.de.